Revisiting CONCATENATE-SEQUENCE

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Abstract

While doing work to support user-extensible sequences (Rhodes, 2007), it was discovered that the ANSI CL standard forbids integration of certain functions with not only user-extensible sequences but also implementation extensions of **sequence**. Irrespective of the future of user-extensible sequences, we argue that the restriction on implementations imposed by the wording adopted is too stringent, and propose an alternative.

1 Introduction

In the X3J13 Issue CONCATENATE-SEQUENCE (Pitman, 1991), the ANSI CL committee worried about various cases of sequence type specifiers passed to the five functions make-sequence, map, merge, concatenate and coerce. The essential problem which the CONCATENATE-SEQUENCE Issue addresses is that a type specifier can specify a recognizable subtype of sequence without unambiguously specifying a concrete sequence type, needed because, except for a special case in coerce, these functions must create an object of the specified type.

For instance, the type sequence itself is a recognizable subtype of sequence; however, the desire was that the call (make-sequence 'sequence 8) should be in error; other such ambiguous types can be constructed, such as (simple-array (*) *), (or bit-vector string), and (and sequence (not (eql "foo"))); although types involving conjunction, disjunction and negation are not required to be recognizable subtypes of sequence, most current implementations recognize these examples as such.

However, the ANSI CL standard also specifies that an implementation may offer subtypes of sequence that are not list and vector:

The types [sic] vector and the type list are disjoint subtypes of type sequence, but are not necessarily an exhaustive partition of sequence.

Pitman and Chapman (1994, System Class sequence)

Historically, this does not appear to have been a popular field for implementation extension; at the time of writing, the author knows of no implementation purporting to conform to Common Lisp which documents non-standard sequence types, though there exist undocumented hooks in at least GNU CLISP (Haible, 2006) which were used in a pre-CLOS implementation of generalized sequences (Haible, 1988).

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In light of this standard definition of the sequence class, and of the development of userextensible sequences, however, the wording for the Exceptional Situations of make-sequence overreaches the intent of the clarification of the CONCATENATE-SEQUENCE issue:

An error of type type-error must be signaled if the result-type is neither a recognizable subtype of list, nor a recognizable subtype of vector.

Pitman and Chapman (1994, Function make-sequence)

Similar requirements are placed on map, merge, concatenate and coerce.

This requirement does not permit an implementation to extend make-sequence to type designators for non-standard sequences, which does not seem to have been the intent behind the CONCATENATE-SEQUENCE issue. We therefore propose the clarification, presented in the style of an issue in the next section.

2 Issue CONCATENATE-SEQUENCE-AGAIN

Issue: CONCATENATE-SEQUENCE-AGAIN.

References: coerce, concatenate, make-sequence, map, merge, Pitman (1991).

Category: Clarification / Change.

Problem Description: The specification says that an error must be signalled in cases when a type specifier passed to make-sequence is not a recognizable subtype of either list or vector. This prevents integration of non-standard sequence types, expressly permitted by the description of sequence, with the standardized sequence functions.

This also affects coerce, concatenate, map and merge.

Proposal (CONCATENATE-SEQUENCE-AGAIN:GENERALIZE):

- Remove from make-sequence, merge, map and concatenate the requirement that "An error of type type-error must be signaled if the result-type is neither a recognizable subtype of list, nor a recognizable subtype of vector."
- Specify that if a type specifier is a recognizable subtype of sequence, and is recognized by the implementation as specifying a concrete subtype of sequence, then a sequence of the specified type is returned from coerce, concatenate, make-sequence, map and merge, subject to the constraints on the type specifier agreeing with the required length of the result sequence.

Rationale: This allows implementors to make extensions of **sequence**, as seems to have been the original intent.

Test Case: No portable test case.

Current Practice: Effectively compatible with both the standard as specified and this proposal, as no implementation extends **sequence** as of the time of writing.

Cost to Implementors: None.

Cost to Users: Minimal. Users can no longer have the guarantee that code of the form

never causes the assertion to fail.

Cost of Non-Adoption: The specification remains inconsistent.

Benefits: A natural way of providing extensions for the sequence type.

Aesthetics: Minimal.

References

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